Modular Multiplication in the Residue Number System

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ΒY

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Abstract

Public-key cryptography is a mechanism for secret communication between parties who have never before exchanged a secret message. This thesis contributes arithmetic algorithms and hardware architectures for the modular multiplication $Z = A \times B \mod M$. This operation is the basis of many public-key cryptosystems including RSA and Elliptic Curve Cryptography. The *Residue Number System* (RNS) is used to speed up long word length modular multiplication because this number system performs certain long word length operations, such as multiplication and addition, much more efficiently than positional systems.

A survey of current modular multiplication algorithms shows that most work in a positional number system, e.g. binary. A new classification is developed which classes these algorithms as Classical, Sum of Residues, Montgomery or Barrett. Each class of algorithm is analyzed in detail, new developments are described, and the improved algorithms are implemented and compared using FPGA hardware.

Few modular multiplication algorithms for use in the RNS have been published. Most are concerned with short word lengths and are not applicable to public-key cryptosystems that require long word length operations. This thesis sets out the hypothesis that each of the four classes of modular multiplication algorithms possible in positional number systems can also be used for long word length modular multiplication in the RNS; moreover using the RNS in this way will lead to faster implementations than those which restrict themselves to positional number systems. This hypothesis is addressed by developing new Classical, Sum of Residues and Barrett algorithms for modular multiplication in the RNS. Existing Montgomery RNS algorithms are also discussed.

The new Sum of Residues RNS algorithm results in a hardware im-

plementation that is novel in many aspects: a highly parallel structure using short arithmetic operations within the RNS; fully scalable hardware; and the fastest ever FPGA implementation of the 1024-bit RSA cryptosystem at 0.4 ms per decryption.

Publications

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Publications in Submission

- 1. Yinan Kong and Braden Phillips, "Modular Reduction and Scaling in the Residue Number System Using Multiplication by the Inverse", submitted to IEEE Transactions on VLSI Systems in November 2008.
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- Braden Phillips and Yinan Kong, "Highly Parallel Modular Multiplication in the Residue Number System using Sum of Residues Reduction", submitted to Journal of Applicable Algebra in Engineering, Communication and Computing in June 2008.
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Nomenclature

 $\langle X \rangle_M$ The operation X mod M.

- D The dynamic range of a RNS.
- M The modulus of a modular multiplication, typically n bits.
- m_i The *i*th RNS channel modulus.
- N The number of RNS channels.
- n The wordlength of M.
- w The RNS channel width.
- $\lceil X \rceil$ The ceiling of X. The smallest integer greater than or equal to X.
- |X| The floor of X. The largest integer smaller than or equal to X.
- BE Base Extension.
- CRT Chinese Remainder Theorem.
- DSP Digital Signal Processing.
- ECC Elliptic Curve Cryptography.
- LUC Look-Up Cycle.
- LUT Look-Up Table.
- LUT Look-Up Table

- MRS Mixed Radix Number System.
- ${
 m QDS}$ Quotient Digit Selection.
- RNS Residue Number System.
- RSA RSA Cryptography.

Contents

1	Intr	oducti	ion	1	
	1.1	Thesis	Outline	2	
	1.2	Contri	ibution	4	
2	Bac	kgrou	nd	7	
	2.1	Residu	le Number Systems	8	
		2.1.1	RNS Representation	8	
		2.1.2	Conversion between RNS and Positional Number Sys-		
			tems	9	
		2.1.3	RNS Arithmetic	14	
		2.1.4	Moduli Selection	15	
		2.1.5	Base Extension	16	
	2.2	RSA I	Public-Key Cryptography	18	
		2.2.1	Public-Key Cryptography	18	
		2.2.2	The RSA Cryptosystem	19	
		2.2.3	Exponentiation	20	
3	Fou	r Way	rs to Do Modular Multiplication in a Positional		
	Number System 23				
	3.1	Introd	ucing Four Ways to Do Modular Multiplication	26	
		3.1.1	Classical Modular Multiplication	27	
		3.1.2	Sum of Residues Modular Multiplication	29	
		3.1.3	Barrett Modular Multiplication	30	
		3.1.4	Montgomery Modular Multiplication	31	

CONTENTS

	3.2	Reinv	igorating Sum of Residues	33
		3.2.1	Tomlinson's Algorithm	33
		3.2.2	Eliminating the Carry Propagate Adder	35
		3.2.3	Further Enhancements	37
		3.2.4	High Radix	40
		3.2.5	Summary of the Sum of Residues Modular Multiplica-	
			tion	40
	3.3	Bound	ls of Barrett Modular Multiplication	42
		3.3.1	Bound Deduction	42
		3.3.2	Performance at Different Word Lengths	45
	3.4	Montg	gomery Modular Multiplication on FPGA	47
		3.4.1	Separated Montgomery Modular Multiplication Algo-	
			rithm	47
		3.4.2	High Radix Montgomery Algorithm	47
		3.4.3	Bound Deduction	54
		3.4.4	Interleaved vs. Separated Structure	57
		3.4.5	Trivial Quotient Digit Selection	58
		3.4.6	Quotient Digit Pipelining	62
	3.5	Modu	lar Multiplications within RNS Channels	69
4	Fou	r Way	vs to Do Modular Multiplication in the Residue	;
	Nui	mber S	System	71
	4.1	Short	Word Length RNS Modular Multiplications using Look-	
		Up Ta	ables	73
		4.1.1	RNS Modular Reduction, Scaling and Look-Up Tables	74
		4.1.2	Existing RNS Scaling Schemes using Look-Up Tables .	75
		4.1.3	A New RNS Scaling Scheme using Look-Up Tables	88
	4.2	RNS (Classical Modular Multiplication	93
		4.2.1	The Core Function	93
		4.2.2	A Classical Modular Multiplication Algorithm in RNS	
			using the Core Function	97
		4.2.3	Examples	105

	4.3	RNS	Sum of Residues Modular Multiplication 109			
		4.3.1	Sum of Residues Reduction in the RNS			
		4.3.2	Approximation of α			
		4.3.3	Bound Deduction			
		4.3.4	The RNS Sum of Residues Modular Multiplication Al-			
			gorithm			
	4.4	RNS	Barrett Modular Multiplication			
		4.4.1	RNS Modular Reduction using Barrett Algorithm 121			
		4.4.2	The Algorithm			
	4.5	RNS	Montgomery Modular Multiplication			
		4.5.1	Montgomery Modular Reduction in RNS			
		4.5.2	The Algorithm $\ldots \ldots 136$			
		4.5.3	A Comparison between RNS Barrett and Montgomery			
			Modular Multiplication			
5	Implementation of RNS Sum of Residues Modular Multipli-					
	cati		- 141			
	5.1	A Sca	alable Structure for Sum of Residues Modular Multipli-			
		catior	n in RNS			
		5.1.1	A 4-Channel Architecture			
		5.1.2	The Scaled Architecture for Modular Multiplication 144			
	5.2	Imple	mentation Results			
6	Conclusion and Future Perspectives 153					
	6.1	Concl	usion			
	6.2	Futur	e Perspectives			
Bi	ibliog	graphy	157			

List of Figures

2.1	Conversion from the RNS to the Mixed Radix Number System (MRS)	12
3.1	An Architecture for Tomlinson's Modular Multiplication Al- gorithm [Tomlinson89]	34
3.2	Modified Sum of Residues Modular Multiplier Architecture	36
3.3	<i>n</i> -bit Carry Save Adders	37
3.4	New Sum of Residues Modular Multiplier Architecture	38
3.5	An Example for New Sum of Residues Modular Multiplication	
	$n = 4, A = (1111)_2, B = (1011)_2$ and $M = (1001)_2$	39
3.6	New High-Radix Sum of Residues Modular Multiplier Archi-	
	tecture \ldots	41
3.7	Delays and Areas of Improved Barrett Modular Multiplication	
	Algorithm	46
3.8	Delays of Montgomery Modular Multiplication at Different	
	Radices	54
3.9	Delays and Areas of Interleaved and Separated Montgomery	
	Modular Multiplication at Different Radices	59
3.10	Shortest Delays of Separated Montgomery Multiplication Al- gorithm using Trivial QDS at n from 12 to 32	62
3.11	Shortest Delays of Separated Montgomery Multiplication Al-	
	gorithm using Quotient Digit Pipelining at n from 13 to 32	67

LIST OF FIGURES

•	3.12	Shortest Delays of Separated Montgomery Multiplication Al- gorithm with & without Quotient Digit Pipelining at n from	
			68
•	3.13	Delays and Areas of Four Classes of Modular Multipliers in	
		Binary	70
2	4.1	RNS Scaling using Look-Up Tables	75
2	4.2	The 3 Modulus RNS Scaler	77
2	4.3	RNS Scaling using Estimation for Channels $S+1 \leq i \leq N$	80
2	4.4	RNS Scaling using a Decomposition of the CRT (two redun-	
		dant channels with channel m_i)	83
2	4.5	RNS Scaling using Parallel Base Extension (BE) Blocks	85
2	4.6	Parallel Architecture to Perform $Y = \lfloor \frac{X}{M} \rfloor \approx \sum_{i=1}^{N} f(x_i)$	86
2	4.7	One Channel of a Parallel Residue Arithmetic Process using	
		Memories with Addressing Capacity $= r. \ldots \ldots \ldots$	87
2	4.8	A New Architecture to Perform Short Word length Scaling in	
		RNS	90
2	4.9	A Typical Core Function C(X): $D=30030;C(D)=165$	95
2	4.10	The h Least Significant Bits to be Maintained during the Mod-	
		ular Multiplication $\langle x_i C(\sigma_i) \rangle_{2^h}$	02
2	4.11	An End-Around Adder to Compute $\langle \theta_i + \theta_{i+1} \rangle_{2^k-1}$	26
2	4.12	Architecture for Scaling by 2 in RNS	28
2	4.13	High-Radix Architecture for Scaling by 2^l in RNS	29
2	4.14	Delay of Channel-Width Multiplication and Addition modulo	
		a General Modulus and Multiplication and Addition modulo	
		$2^k - 1$ used for RNS Barrett Algorithm against RNS Channel	
		Width w	32
2	4.15	Delay of Barrett and Montgomery RNS Modular Multiplica-	
		tion against RNS Channel Width w	38
ļ	5.1	A 4-Channel RNS Modular Multiplier	45

5.2	A 4-Channel Modular Multiplier Block Suitable for a Scalable
	Array of Blocks
5.3	A 12-Channel Modular Multiplier Built from a 3×3 Array of
	4-Channel Blocks

List of Tables

2.1	The Pre-Computed Constants of RNS $\{3, 5, 7\}$	9
2.2	An Example of Right-to-Left Modular Exponentiation	23
3.1	Montgomery Modular Reduction Process at Radix 2	49
3.2	Montgomery Modular Reduction Process at Radix 4	50
3.3	Montgomery Modular Reduction Process at Radix 8 $(l = \left\lceil \frac{n}{k} \right\rceil)$	52
3.4	Montgomery Modular Reduction Process at Radix 8 $(l = \lfloor \frac{n}{k} \rfloor)$	53
3.5	The k-boundary of Montgomery Reduction for $n \mbox{ from } 12 \mbox{ to } 32$	61
3.6	Shortest Delays of Separated Montgomery Multiplication Al- gorithm using Trivial QDS at n from 12 to 32	61
3.7	k and d in Quotient Digit Pipelining of Montgomery Algorithm	
	for n from 12 to 25 \ldots	66
3.8	Shortest Delays of Separated Montgomery Multiplication Algorithm using Quotient Digit Pipelining at $n \mbox{ from 13 to 32}$.	67
4.1	A Comparison of 6 Short Word length RNS Scaling Schemes using LUTs	92
4.2	An Example of Pre-Computed Parameters for the Classical	
	Modular Reduction Algorithm in RNS using the Core Function	106
4.3	An Example of the Classical Modular Reduction Algorithm in	
	RNS using the Core Function for a case in which $\lfloor Mag(X) \rfloor \ge 1$	107
4.4	An Example of the Classical Modular Reduction Algorithm in	
	RNS using the Core Function for a case in which $\lfloor Mag(X) \rfloor < 1$	108

4.5	The Maximum Possible N against w in RNS Sum of Residues
	Reduction
5.1	The Pre-Computed Constants for the RNS Sum of Residues
	Modular Multiplication Algorithm in Algorithm 4.2 144